

On Levels in Arrangements of Curves, III:

Further Improvements

Timothy Chan

School of CS

U of Waterloo

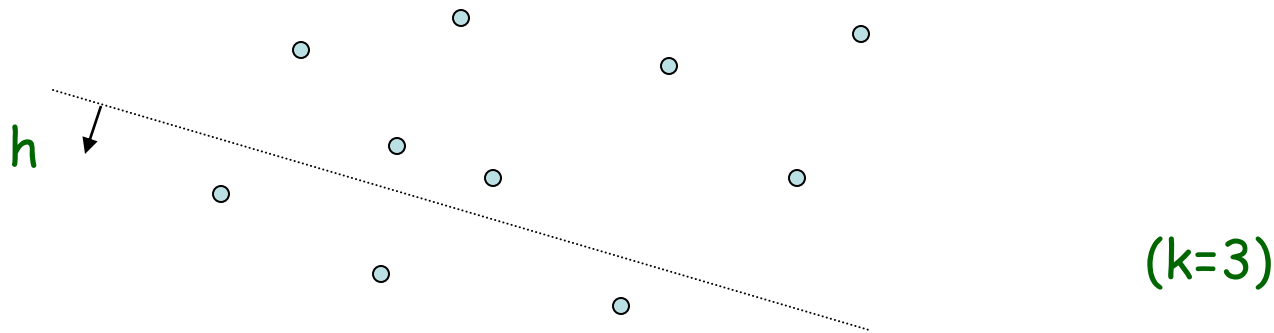
Back Story

- [FOCS'00] "On Levels in Arrangements of Curves"
- [FOCS'03] "On Levels in Arrangements of Curves, II"

- [SODA'05] "... of Surfaces in 3D"
- [SODA'08] "On the Bichromatic k -Set Problem"

The k-Set Problem

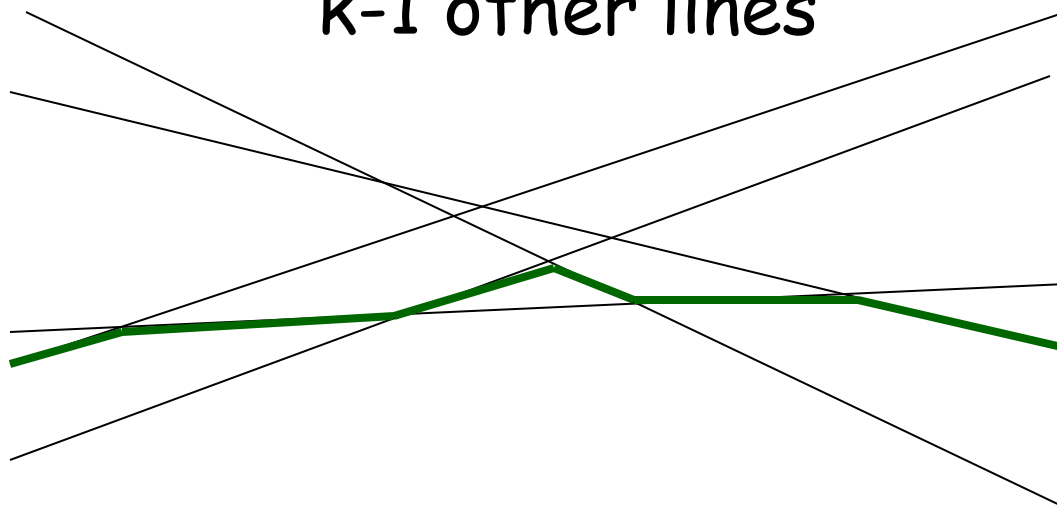
- Given n pts P in 2D & k , a **k-set** is a subset $P \cap h$ of size k for some halfplane h



- How many k -sets possible (as fn of n)?

Equiv. Dual Problem: k -Level of Lines

- Given n lines in 2D & k ,
 k -level := all pts that lie on 1 line & above
 $k-1$ other lines



- How many vertices on k -level?

Known Upper Bds

- Lovász'71 $O(n^{3/2})$
- Gusfield'79 $O(n^{3/2})$
- Edelsbrunner, Welzl'85 $O(n^{3/2})$
- Pach, Steiger, Szemerédi'89 $O(n^{3/2} / \log^* n)$
- Dey'97 $O(n^{4/3})$

(by counting crossings in a geometric graph in primal space)

Known Lower Bds

- Erdős, Lovász, Simmons, Straus'73 $\Omega(n \log n)$
- Klawe, Paterson, Pippenger'82 $n^{2^{\Omega(\sqrt{\log n})}}$
(for k -level of pseudo-lines)
- Tóth'00 $n^{2^{\Omega(\sqrt{\log n})}}$
- Conjecture by Erdős et al.'73:
 $o(n^{1+\varepsilon})$??

Going "Retro"

- Chan [FOCS'03]
 - a completely different pf of $O(n^{3/2})$
(by a simple inequality in dual space)

Going "Retro"

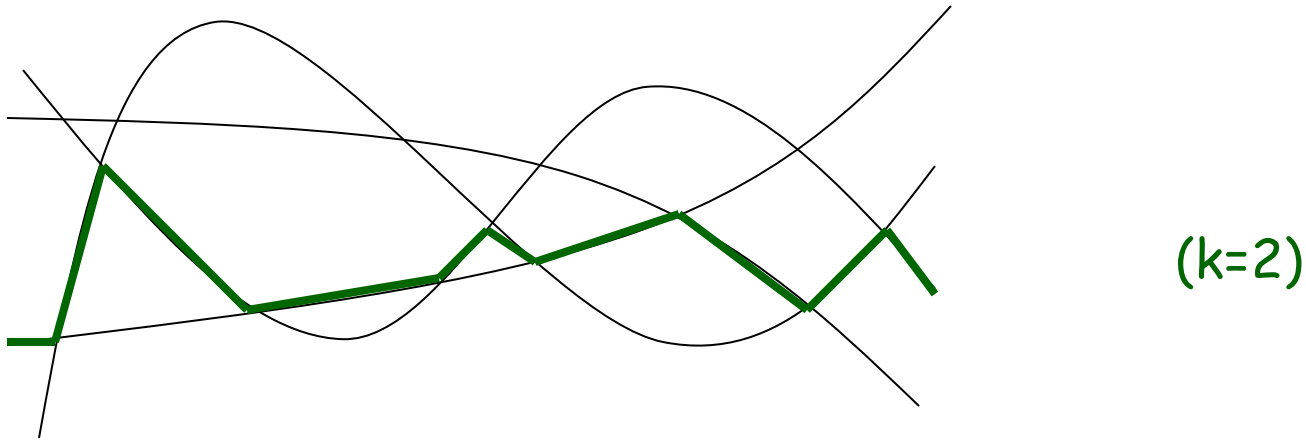
- Chan [FOCS'03]
 - a completely different pf of $O(n^{3/2})$
(by a simple inequality in dual space)
- Today:
 - a new pf of $O(n^{3/2 - \delta})$ for a concrete but small $\delta > 0$
(by refining the inequality)

Why Interesting, Despite of Dey?

- Because the inequality approach
 - is "completely different"
 - is the **only** approach known for k-level of general s-intersecting curves [FOCS'03]
 - gives **current best** bds for k-level of specific curve families [FOCS'03] & for related problems (e.g., k-level of surfaces in 3D [SODA'05], "bichromatic" k-sets in 2D & 3D [SODA'08])

The k -Level Problem for Curves

- Given n x -monotone curves in 2D & k ,
 k -level := all pts that lie on 1 curve &
above $k-1$ other curves



- How many vertices on k -level?

Known Upper Bds

- Pseudo-lines:

Dey + Tamaki, Tokuyama'97 $O(n^{4/3})$

- Pseudo-parabolas:

Tamaki, Tokuyama'95 $O(n^{23/12})$

Dey'97 $O(n^{17/9})$

[FOCS'00] $O(n^{16/9} \log^{2/3} n)$

Agarwal, Nevo, Pach, Pinchasi,
Sharir, Smorodinsky'02 $O(n^{26/15} \log^{2/3} n)$

[FOCS'03] $O(n^{8/5})$

Marcus, Tardos'04 $O(n^{3/2} \log^2 n)$

Today $O(n^{3/2} \log n)$

Known Upper Bds (cont'd)

- Pseudo-segments:

[FOCS'00]

$$O(n^{4/3} \log^{2/3} n)$$

Today

$$O(n^{4/3} \log^{1/3-\delta} n)$$

- Degree- s polynomials:

[FOCS'00]

$$O(n^{2 - O(1/2^s)})$$

- General s -intersecting curves/curve segments:

[FOCS'03]

$$O(n^{2 - 1/(2s)}) \text{ for } s \text{ odd}$$

$$O(n^{2 - 1/(2s-2)}) \text{ for } s \geq 4 \text{ even}$$

Today

$$O(n^{2 - 1/(2s) - \delta_s}) \text{ for } s \text{ odd}$$

$$O(n^{2 - 1/(2s-2) - \delta_s}) \text{ for } s \geq 4 \text{ even}$$

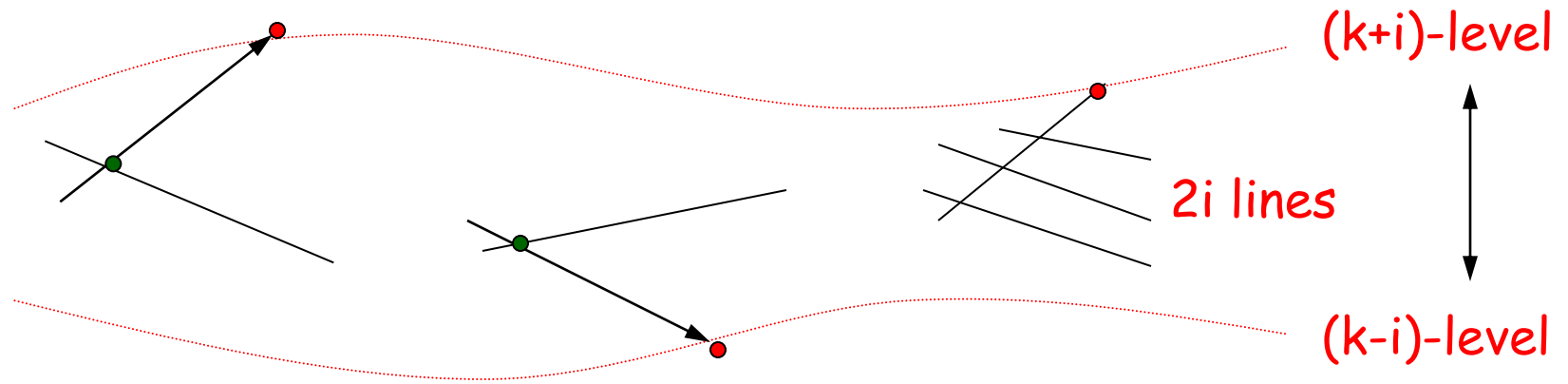
The Inequality Approach [FOCS'03]: The (Pseudo-)Line Case

- Let $t_i = \#$ vertices at levels in interval $(k-i, k+i)$
- Let $\Delta t_i = t_{i+1} - t_i$

• Obs:

$$t_i \lesssim 2i \Delta t_i$$

• Pf:



The Inequality Approach (cont'd)

- Solve $t_i \lesssim 2i(t_{i+1} - t_i)$

$$\Rightarrow t_i \lesssim [2i/(2i+1)] t_{i+1}$$

$$\Rightarrow t_1 \lesssim 2/3 \cdot 4/5 \cdot 6/7 \cdots [2n/(2n+1)] t_n$$

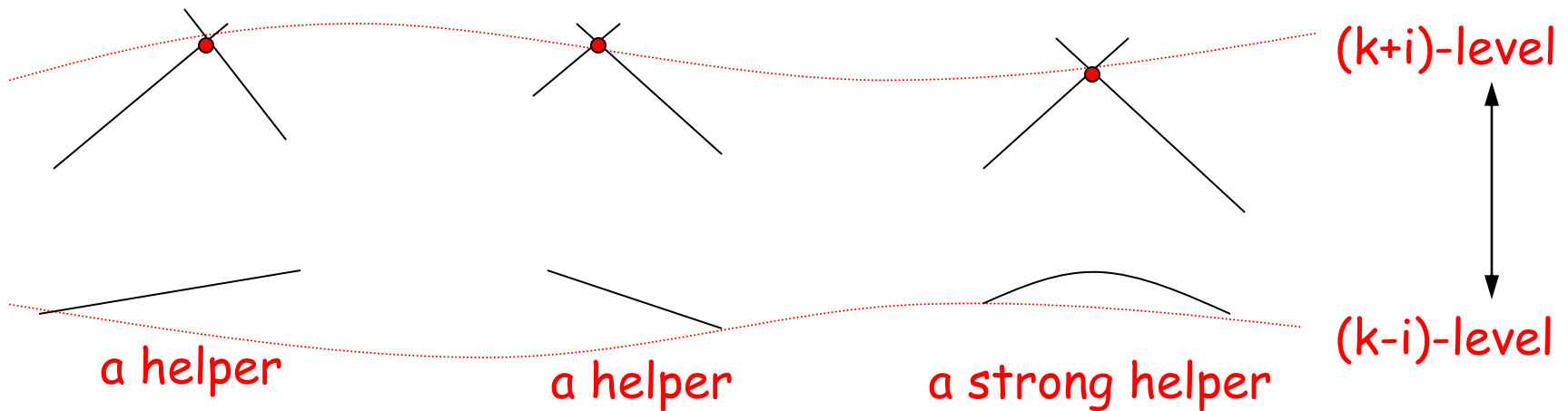
$$\leq O(1/n^{1/2}) t_n$$

$$= \boxed{O(n^{3/2})}$$

- Rmk: $O(n^{3/2} + B)$ for "almost" pseudo-lines/segments with B "bad" pairs

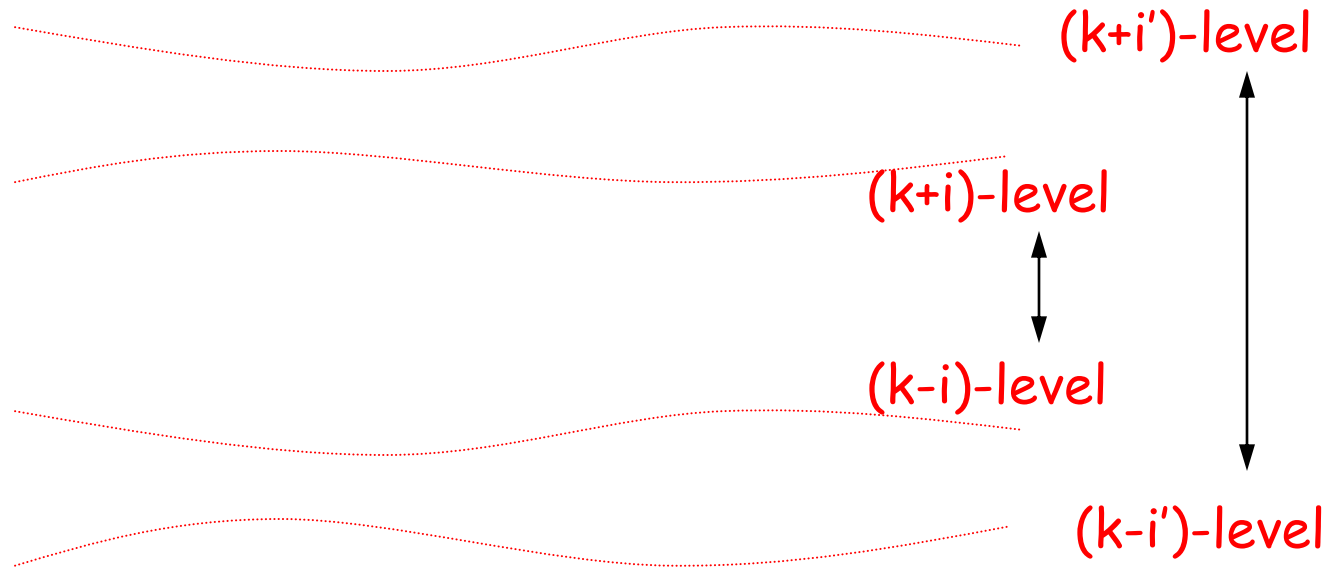
New Idea 1: Look for Slack

- Define **helpers**:

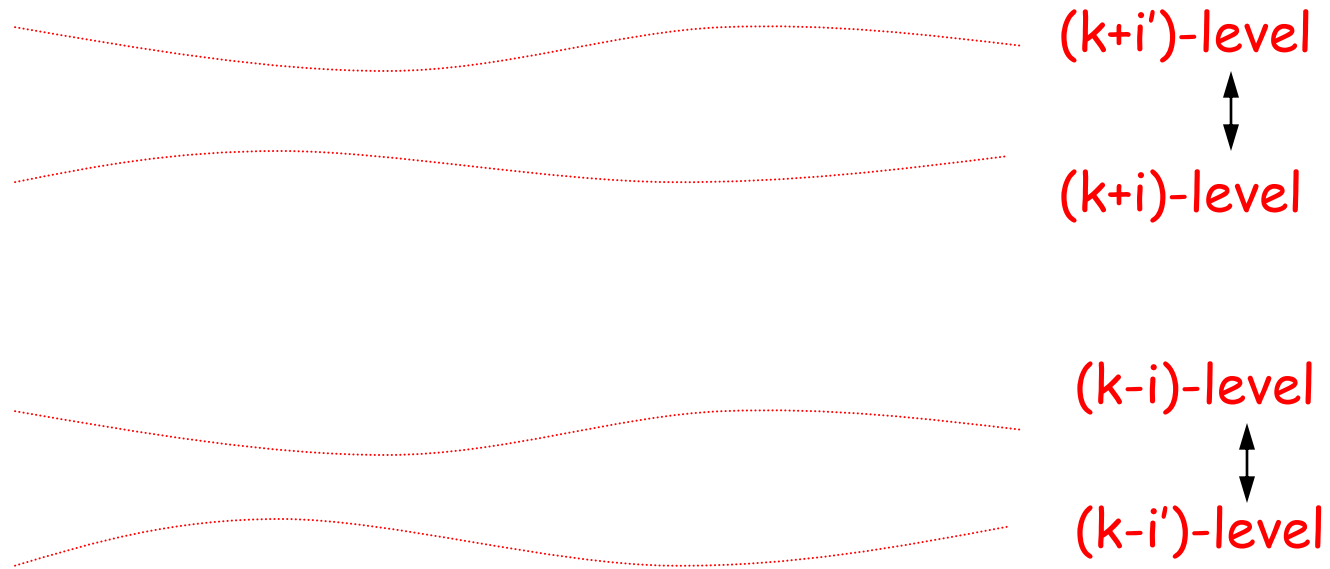


- **Obs:** $t_i \lesssim 2i \Delta t_i - \Omega(\# \text{ helpers})$
- But how to find helpers?

New Idea 2: Look at Two Intervals



New Idea 3: Look at The "Side" Intervals

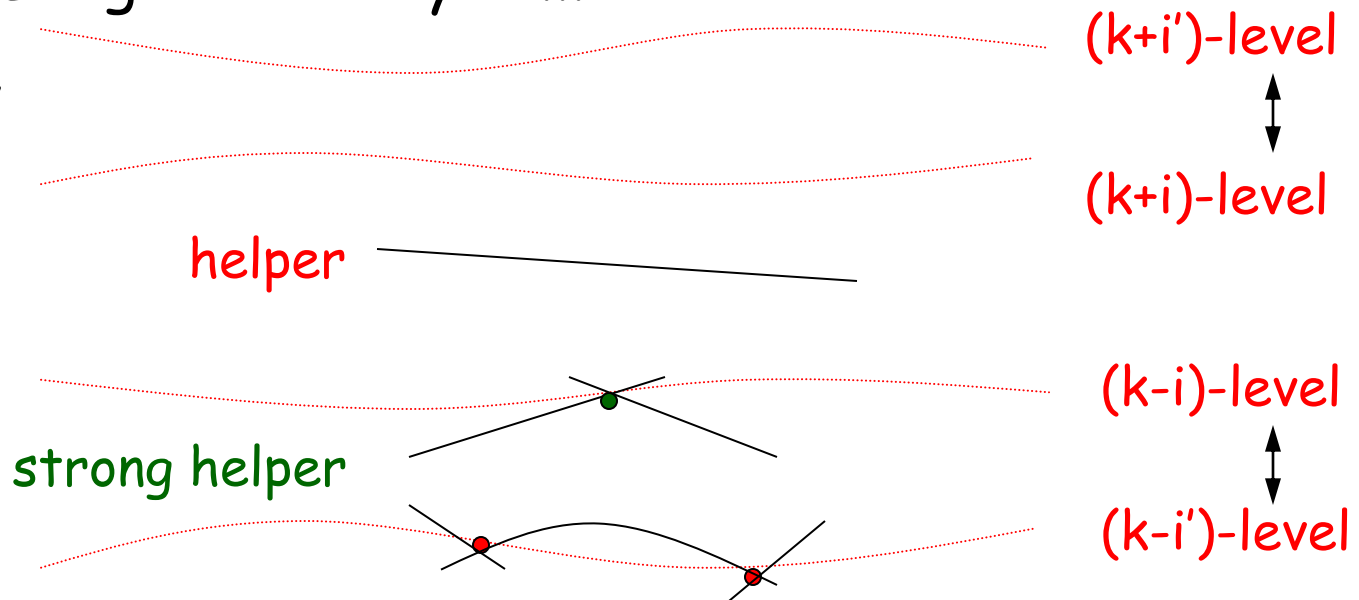


Obs: $t_{i'} - t_i \gtrsim (i' - i) / 2 (\Delta t_i + \Delta t_{i'})$
- $O(\# \text{ strong helpers in side intervals})$

New Idea 4: Surprise!

- Lemma: # helpers in the two intervals
 $\gtrsim c \cdot$ # strong helpers in side intervals
- Pf: Long case analysis...

e.g.,



New Inequality

$$t_i + t_{i'} \lesssim 2i \Delta t_i + 2i \Delta t_{i'} - c [(i'-i)/2 (\Delta t_i + \Delta t_{i'}) - (t_{i'} - t_i)]$$

• Solve

$$\Rightarrow t_i \leq O(1/n^{1/2 + \delta}) t_n = O(n^{3/2 - \delta})$$

Further Results/Open Questions

- Another, simpler way to refine the inequality for s -intersecting curves for odd $s \geq 3$ (but not curve segments or even s)
- Original k -set problem: $O(n^{4/3 - \delta})$??
- Another observation:
 - Better bds for const $n \Rightarrow$ better asymptotic upper bds
- k -level of curves: better upper or lower bds ??